# An Experimental Comparison of Partitioning Strategies in Distributed Graph Processing

26-Feb-2019

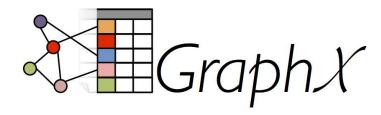
Presented by: Noshin Nawar Sadat



#### **Graph Processing Systems**















#### **Graph Partitioning**

Graph processing systems run in distributed manner to handle large graphs.

To distribute computations, partition graphs by assigning graph edges or vertices to individual machines.



#### The Problem

Plethora of graph processing systems.

Each offer their own graph partitioning strategies.

Even after user chooses a system, there are different partitioning strategies to choose from.



#### **The Problem**

Choosing the right partitioning strategy for distributed graph processing systems



#### Why is the problem important?

Choice of partitioning strategy affects

- Performance
  - Replication factor
  - Completion time (Ingress/Partitioning time, Computation time)
- Resource usage
  - Memory usage
  - Network I/O



#### **Paper Contributions**

- Experimentally compared partitioning strategies of three popular distributed graph processing (vertex-cut based) systems:
  - PowerGraph
  - PowerLyra
  - GraphX
- Presented rules of thumb to help pick partitioning strategies within each system for different use cases.
- Implemented PowerGraph and GraphX's strategies into PowerLyra



## **Experimental Methodology: Clusters**

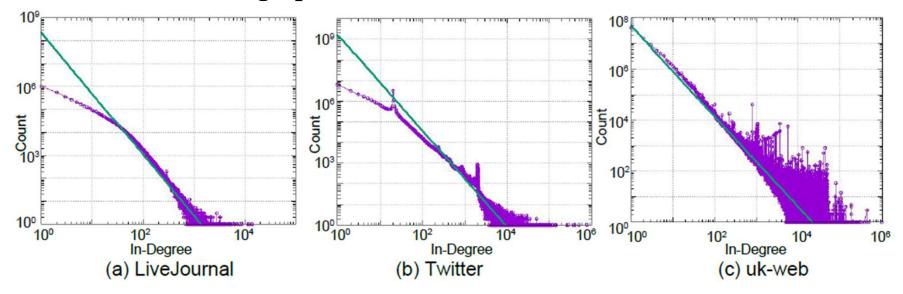
- PowerGraph and PowerLyra
  - A local cluster of 9 machines
  - EC2 consisting of 16 m4.2xlarge instances
  - EC2 consisting of 25/m4.2xlarge instances
- GraphX
  - Local cluster of 10 machines

Grid requires number of partitions to be a perfect square.



# **Experimental Methodology: Datasets**

- Consisted of:
  - Low degree graphs (road-net-CA, road-net-USA)
  - Heavy tailed graphs (LiveJournal, Enwiki-2013, Twitter)
  - Power-law like graphs (UK-web)



# **Experimental Methodology: Metrics**

- Ingress time
- Computation time
- Replication factor
- System wide resource usage
  - memory consumption
  - CPU utilization
  - network I/O usage



## **Experimental Methodology: Applications**

- PageRank
- Weakly connected components
- K-core decomposition
- Single Source Shortest Path (SSSP)
- Simple coloring



- Random
- Oblivious
- Constrained
  - Grid
  - Perfect Difference Sets (PDS)
- High Degree replicated
   First (HDRF)



- Random
- Oblivious
- Constrained
  - Grid
  - Perfect Difference Sets (PDS)
- High Degree replicated
   First (HDRF)

- Edge's hash
  - = function of its vertices
- Fast
- Parallelizable
- Even edge distribution
- Creates large number of mirrors



- Random
- Oblivious
- Constrained
  - Grid
  - Perfect Difference Sets (PDS)
- High Degree replicated
   First (HDRF)

- Based on a greedy heuristic
- Goal : Keep replication factor low
- Incrementally and greedily place edges
- Requires some information about previous assignments
- Not a trivial strategy to parallelize and distribute

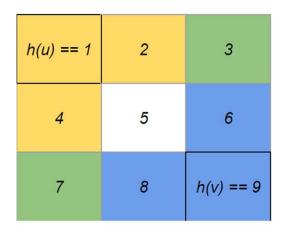


- Random
- Oblivious
- Constrained
  - Grid
  - Perfect Difference Sets (PDS)
- High Degree replicated
   First (HDRF)

- Hash edges
- Restrict edge placement based on vertex adjacency:
  - Constraint set of vertexv = S(v)
  - Place edge (u,v) in partition belonging to set S(u)∩S(v)



- Random
- Oblivious
- Constrained
  - Grid
  - Perfect Difference Sets (PDS)
- High Degree replicated
   First (HDRF)



- S(v) = all machines in the row and column of the machine that v hashed to
- Upper bound for replication factor:  $2\sqrt{N-1}$



- Random
- Oblivious
- Constrained
  - Grid
  - Perfect Difference Sets (PDS)
- High Degree replicated First (HDRF)

- Uses PDS to generate constraint sets
- Requires  $(p^2 + p + 1)$  machines, where p is prime.

Couldn't meet requirement for number of machines of both Grid and PDS



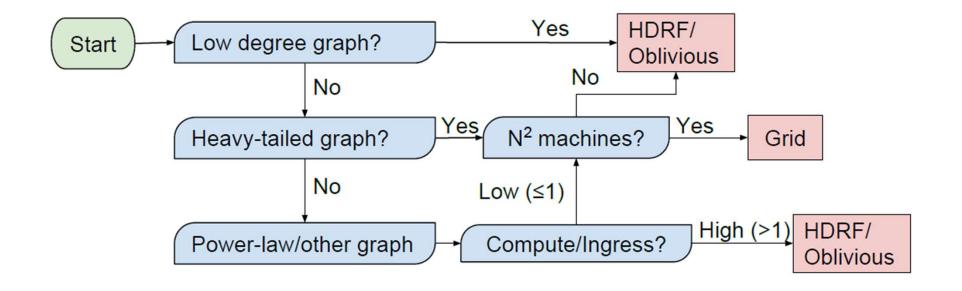
- Random
- Oblivious
- Constrained
  - Grid
  - Perfect Difference Sets (PDS)
- High Degree replicated First (HDRF)

- Similar to Oblivious
- Assigns edges by looking at both partition size and vertex degrees
- Prefers replicating high degree vertices



## **PowerGraph: Results**

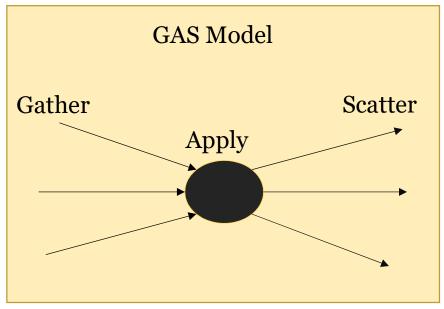
| Strategy | PageRank (Conv.) |         |       | K-Core Decomp. |         |        |
|----------|------------------|---------|-------|----------------|---------|--------|
|          | ingress          | compute | total | ingress        | compute | total  |
| Grid     | 206.4            | 146.0   | 352.4 | 203.6          | 3794.9  | 3998.5 |
| HDRF     | 322.0            | 103.6   | 425.6 | 320.6          | 3225.1  | 3545.7 |





- Random
- Oblivious
- Constrained
  - Grid
  - Perfect Difference Sets (PDS)
- Hybrid
- Hybrid-Ginger





- Hybrid
- Hybrid-Ginger

- Low degree destinations:
   hash destination vertex
  - Colocate with all in-edges
- High degree destinations : hash source vertex
  - Colocate with all out-edges



- Random
- Oblivious
- Constrained
  - Grid
  - Perfect Difference Sets (PDS)
- Hybrid
- Hybrid-Ginger

- Two phases:
  - 1. Performs edge cuts on all vertices and updates degree counters
  - 2. Reassignment phase: performs vertex cuts on the vertices whose degree is above a certain threshold



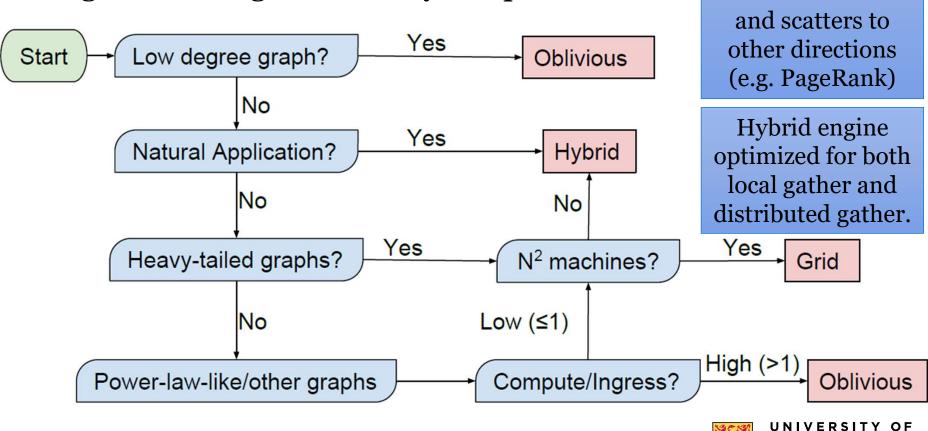
- Random
- Oblivious
- Constrained
  - Grid
  - Perfect Difference Sets (PDS)
- Hybrid
- Hybrid-Ginger

- A greedy streaming edgecut strategy
- Three phases:
  - 1 & 2. Same as Hybrid
  - 3. Put low degree vertex v in partition which has more of its in-neighbors
  - Gather locally



#### **PowerLyra: Results**

 Hybrid-Ginger consistently showed slower ingress and higher memory footprint.



PAGE 15

**Natural** 

**Applications:** 

Vertex gathers

from one direction

- Random
- Canonical Random
- 1D Edge Partitioning
- 2D Edge Partitioning



- Random
- Canonical Random
- 1D Edge Partitioning
- 2D Edge Partitioning

- Hashing the source and vertex IDs to assign edges
- Edges (u,v) and (v,u) go to the same partition



- Random
- Canonical Random
- 1D Edge Partitioning
- 2D Edge Partitioning

- Hashing the source and vertex IDs in canonical direction to assign edges
- Edges (u,v) and (v,u) do not necessarily go to same partition
- Similar to Random of PowerGraph/PowerLyra



- Random
- Canonical Random
- 1D Edge Partitioning
- 2D Edge Partitioning

- Hash all edges by their source vertex
- Ensure all edges with the same source are in same partition
- Similar to PowerLyra's
   Hybrid's low degree vertex partitioning



- Random
- Canonical Random
- 1D Edge Partitioning
- 2D Edge Partitioning

- Arrange all partitions into square matrix
- Pick column on the basis of source vertex's hash and the row on the basis of destination vertex's hash
- Ideal if number of partitions = perfect square
- Similar to PowerGraph/PowerLyra's Grid



#### **GraphX**: Results

- Canonical random for low degree and high diameter graphs (e.g. road networks)
- 2D Edge Partitioning for power-law like graphs (e.g. Enwiki-2013)



#### **PowerLyra: All Strategies**

- HDRF
- 1D Edge partitioning
- 2D Edge partitioning
- Asymmetric Random (GraphX's Random)
- 1D-Target (new strategy)
  - Hash edges by destination vertex and colocate them



#### **PowerLyra: All Strategies - Results**

- No change in PowerLyra's decision tree
  - HDRF same as Oblivious
- Both Asymmetric Random and Random's performance were the worst
- PowerLyra's Hybrid engine enhances 2D edge partitioning



#### **Summary**

- No best fit for all situations
- Choice depends on:
  - Degree of distributed graph
  - Type and duration of application
  - Cluster size
- Choice significantly affects resource usage and application runtime
- Partitioning strategies tightly integrated with underlying engine perform better.

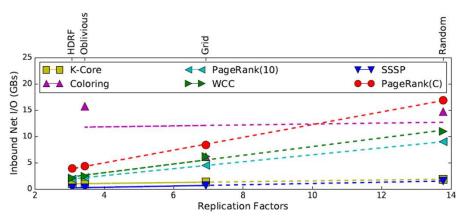


#### **Discussion**

- Why did they choose PowerGraph, PowerLyra and GraphX?
- Was leaving out the PDS partitioning strategy a good idea?
- PowerGraph can be used with both synchronous and asynchronous engines. In the paper, they mentioned that they used asynchronous engine for simple coloring, which led to longer computation time and higher network I/O usage. Why did they not use synchronous engine for it?
- Why did Hybrid-Ginger give such bad performance?



#### **Discussion**



**Figure 3:** Incoming Network IO vs. Replication Factors. (PowerGraph, EC2-25, UK-Web).

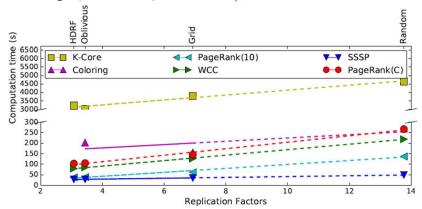


Figure 4: Computation Time in seconds vs. Replication Factors. (PowerGraph, EC2-25, UK-Web).

